Software performance modelling using PEPA nets

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ABSTRACT

Modelling and analysing distributed and mobile software systems is a challenging task. PEPA nets—coloured stochastic Petri nets—are a recently introduced modelling formalism which clearly capture important features such as location, synchronisation and message passing. In this paper we describe PEPA nets and the newly-developed platform support for software performance modelling using them. Crucial to this support is the compilation from PEPA nets into Hillston's PEPA stochastic process algebra in order to access the software tools which support the PEPA algebra. In addition to derivation of steady state performance measures, this suite of tools allows properties of the system to be verified using model-checking. We show the application of PEPA nets in the modelling and analysis of a secure Web service.

Categories and Subject Descriptors

D.2.1 [**Software Engineering**]: Requirements / Specifications — *Methodologies* ; D.3.2 [**Programming Languages**]: Language Classifications ; D.2.8 [**Software Engineering**]: Metrics—*performance measures*

General Terms

PEPA nets; performance analysis; mobile objects

1. INTRODUCTION

The correct representation of many modern software systems needs us to distinguish several distinct contexts of computation. These contexts may depend on physical location, operating conditions, or both. For example, in a distributed system interaction which takes place locally may be richer than that available through a remote interface; mobile objects may migrate within a system, significantly altering communication costs; and in fault-tolerant systems the current fault scenario can significantly affect performance.

Design decisions such as the placement of software components on hosts, the separation of client functions from server functions, and movement of code and data from host to host are typically motivated by the need to improve system reliability and performance. Unfortunately, many effective state-of-the-art performance analysis tools do not provide support, for example, for differentiating between simple local communication and the movement or migration of processes which can change the allowable pattern of communication. One formalism which does separate these two mechanisms is the *PEPA nets* modelling language [1].

PEPA nets extend the PEPA stochastic process algebra [2] by allowing PEPA process algebra components to be used as the tokens of a coloured stochastic Petri net. A PEPA component has local state which can be modified by performing timed actions, either individually or in cooperation with another component. These activities define an interface analogous to the interface made up of the methods which can be invoked on an object. Constraints on the nature of the timing information (exponentially distributed random delays) mean that every PEPA model will define a Continuous-Time Markov Chain (CTMC). The addition of the Petri net infrastructure allows the modeller to represent explicitly locations or contexts (the statically-defined places of the net) and movement between them (the dynamically-determined firings of the net). Software components of the system, represented by PEPA components, make up the tokens of the net and move between the places of the net according to a formally defined firing rule.

Furthermore, the semantics of the formalism imposes a restriction on communication based on the structure of the net. Communication is strictly confined within places—only components presently within the same context may interact. Thus interaction between tokens at different places is not allowed; for tokens to interact they must first meet at one of the places of the net. Moreover as a token moves around the net, its possible activities may be altered to reflect its current location. Restricting communication according to these rules provides, for example, a credible formalism for representing mobile code systems where the tokens of a PEPA net model stateful mobile objects under a system of dynamic binding of names.

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As explained above, the places of the PEPA net provide the locations, or contexts of the system. These are themselves modelled by PEPA components, termed *static components*, which, unlike tokens, are unable to move from the place. These immobile parts of the system may represent hardware (such as servers) or software (such as databases). Static components may communicate with tokens while they are present or they can communicate with other static components in the same place. Communication between static components at different places is not allowed.

In this paper we present PEPA nets, which have recently been extended to include more general net structures, and demonstrate their modelling and analysis capabilities on a small example of a mobile agent system and a more substantial case study based on a secure Web service.

There is an extensive suite of tools already supporting the PEPA language. In order to be able to exploit these tools, rather than develop new ones, we have developed an algorithm which compiles a PEPA net model to an equivalent model in PEPA. The difficulty here is transforming a PEPA net model into a model in a language which has no explicit notion of either location or mobility between locations. Despite the complexity of the transformation, we believe it is worthwhile since it provides access to established, trusted tools. Furthermore it is implemented in a software tool and therefore transparent to the user.

Structure of this paper: In the next section we present examples and the definition of the PEPA nets modelling language. This definition extends previous versions, by allowing a more general net structure to be used. We also explain the firing rule, which has been appropriately extended to take account of this more general structure. In Sect. 3 we describe the method of translating PEPA nets to PEPA and present results about the translation. In Sect. 4 we present an example of the algorithm. In Sect. 5 we describe the implementation of our translation method as an extension to the PEPA Workbench for PEPA nets which connects to the tools of the PEPA platform. The objective of the algorithm is to allow PEPA net users to be able to exploit the existing tool support for PEPA. In Sect. 6 we give a brief overview of that tool support. In Sect. 7 we describe our Web Services case study. Conclusions are presented in Sect. 8.

2. PEPA NETS

In this section we provide a brief overview of PEPA nets and the PEPA stochastic process algebra. A fuller description, together with supporting theory and proofs is available in [1] and [2]. The purpose of this summary is to provide enough information about the modelling language to make the present paper self-contained.

The tokens of a PEPA net are terms of the PEPA stochastic process algebra which define the behaviour of components via the activities they undertake and the interactions between them. One example of a PEPA component would be a *File* object which can be opened for reading or writing, have data read (or written) and closed. Such an object would understand the methods *openRead*(), *openWrite*(), *read*(), *write*() and *close*().

A PEPA model shows the order in which such methods can be invoked.

$$
\begin{aligned}
\mathit{File} &\stackrel{\mathit{def}}{=} (\mathit{openRead}, r_o).\mathit{InStream} \\
&\quad + (\mathit{openWrite}, r_o).\mathit{OutStream} \\
\mathit{InStream} &\stackrel{\mathit{def}}{=} (\mathit{read}, r_r).\mathit{InStream} \\
&\quad + (\mathit{close}, r_c).\mathit{File} \\
\mathit{OutStream} &\stackrel{\mathit{def}}{=} (\mathit{write}, r_w).\mathit{OutStream} \\
&\quad + (\mathit{close}, r_c).\mathit{File}
\end{aligned}
$$

Every activity in the model incurs an execution cost which is quantified by an estimate of the (exponentially-distributed) rate at which it can occur (r_o, r_r, r_w, r_c).

Such a description documents a high-level protocol for using *File* objects, from which it is possible to derive properties such as "it is not possible to write to a closed file" and "read and write operations cannot be interleaved: the file must be closed and re-opened first".

A PEPA net is made up of PEPA *contexts*, one at each place in the net. A context consists of a number of *static* components (possibly zero) and a number of *cells* (at least one). Like a memory location in an imperative program, a cell is a storage area to be filled by a datum of a particular type. In particular in a PEPA net, a cell is a storage area dedicated to storing a PEPA component, such as the *File* object described above. The components which fill cells can circulate as the tokens of the net. In contrast, the static components cannot move. A typical place might be the following:

$$\mathit{File}[_] \underset{L}{\bowtie} \mathit{FileReader}$$

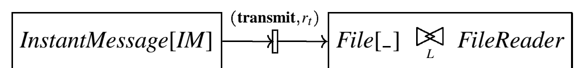
where the *synchronisation set* L in this case is $\vec{\mathcal{A}}(\mathit{File})$, the *complete action type set* of the component, (*openRead*, *openWrite*, ...). This place has a *File*-type cell and a static component, *FileReader*, which can process the file when it arrives.

A PEPA net differentiates between two types of change of state. We refer to these as *firings* of the net and *transitions* of PEPA components. Each are special cases of PEPA activities. Transitions of PEPA components will typically be used to model small-scale (or *local*) changes of state as components undertake activities. Firings of the net will typically be used to model macro-step (or *global*) changes of state such as context switches, breakdowns and repairs, one thread yielding to another, or a mobile software agent moving from one network host to another. The set of all firings is denoted by $\mathcal{A}_f$. The set of all transitions is denoted by $\mathcal{A}_t$. We distinguish firings syntactically by printing their names in boldface.

Continuing our example, we introduce an instant message as a type of transmissible file.

$$\mathit{InstantMessage} \stackrel{\mathit{def}}{=} (\mathbf{transmit}, r_t).\mathit{File}$$

Part of a definition of a PEPA net which models the passage of instant messages is shown below.



An instant message *IM* can be moved from the input place on the left to the output place on the right by the **transmit** firing. In doing so it changes state to evolve to a *File* derivative, which can be read by the *FileReader*.

The syntax of PEPA nets is given in Figure 1. In that grammar S denotes a *sequential component* and P denotes a *concurrent component* which executes in parallel. I stands for a constant which denotes either a sequential or a concurrent component, as bound by a definition.

$$
\begin{array}{rcll}
N & ::= & D^{+}M & \text{(net)} \\
M & ::= & (M_{\mathbf{P}}, \ldots) & \text{(marking)} \\
M_{\mathbf{P}} & ::= & \mathbf{P}[C, \ldots] & \text{(place marking)} \\
D & ::= & I \stackrel{\mathit{def}}{=} S & \text{(component defn)} \\
 & | & \mathbf{P}[C] \stackrel{\mathit{def}}{=} P[C] & \text{(place defn)} \\
 & | & \mathbf{P}[C, \ldots] \stackrel{\mathit{def}}{=} P[C] \underset{L}{\bowtie} P & \text{(place defn)} \\
P & ::= & P \underset{L}{\bowtie} P & \text{(cooperation)} \\
 & | & P/L & \text{(hiding)} \\
 & | & P[C] & \text{(cell)} \\
 & | & I & \text{(identifier)} \\
C & ::= & \text{`_'} & \text{(empty)} \\
 & | & S & \text{(full)} \\
S & ::= & (\alpha, r).S & \text{(prefix)} \\
 & | & S + S & \text{(choice)} \\
 & | & I & \text{(identifier)}
\end{array}
$$

Figure 1: The syntax of PEPA nets

DEFINITION 1 (PEPA NET). *A PEPA net $\mathcal{N}$ is a tuple $\mathcal{N} = (\mathcal{P}, \mathcal{T}, I, O, \ell, \pi, \mathcal{C}, D, M_0)$ such that*

- *$\mathcal{P}$ is a finite set of places;*
- *$\mathcal{T}$ is a finite set of net transitions;*
- *$I : \mathcal{T} \to \mathcal{P}$ is the input function;*
- *$O : \mathcal{T} \to \mathcal{P}$ is the output function;*
- *$\ell : \mathcal{T} \to (\mathcal{A}_f, \mathbb{R}^+ \cup \{\top\})$ is the labelling function, which assigns a PEPA activity ((type, rate) pair) to each transition. The rate determines the negative exponential distribution governing the delay associated with the transition;*
- *$\pi : \mathcal{A}_f \to \mathbb{N}$ is the priority function which assigns priorities (represented by natural numbers) to firing action types;*
- *$\mathcal{C} : \mathcal{P} \to P$ is the place definition function which assigns a PEPA context, containing at least one cell, to each place;*
- *D is the set of token component definitions;*
- *M_0 is the initial marking of the net.*

We use the notation

$$P_1 \xrightarrow{(\alpha, r)} P_2$$

to capture the information that there is a transition connecting place P_1 to place P_2 labelled by (α, r).

The structured operational semantics, given in [1], gives a precise definition of the possible evolution of a PEPA net, and show how a CTMC can be derived, treating each marking as a distinct state.

We define the firing rule of PEPA nets to respect the net structure in the usual way (one token from each input arc, one token to each output arc) but also to take into consideration the ability of tokens to participate in the firing (can they perform an activity of the correct type?), and the availability of vacant cells of the appropriate type in the output places. Note that in previously published papers on PEPA nets the net structure was restricted to form a state machine, thus removing the possibility of more than one token moving at once. We have now removed this restriction but still require that the net is *balanced* in the sense that, for each transition, the number of input arcs is equal to the number of output arcs. In classical Petri nets tokens are identitiless, and can be viewed as being consumed from input places and created into output places for each firing. In contrast, in PEPA nets our tokens have state and identity, and we view them as *passing through* net-level transitions. For each firing there must be as many output tokens as there were input tokens. This generalisation has increased the expressiveness of the formalism.

DEFINITION 2 (ENABLING). *An enabling is a mapping of places to tokens. A transition t has an* enabling *of firing type α, $\mathrm{E}(t, \alpha)$, if for each input place $\mathbf{P}_i$ of t there is a token T in the current marking of $\mathbf{P}_i$, which has a one-step α-derivative, T'.*

Note that there may be several enablings for a given transition firing in any particular marking, as the enabling selects one token to fire from each input place[1], and there may be more than one eligible token at each input place.

Since it is important that each fired token has a vacant cell to go into after the firing, we define a corresponding notion of *output*. A transition has an output if, in the current marking, there is at least one vacant cell in each output place.

DEFINITION 3 (OUTPUT). *For any transition t, an* output, *denoted $O(t)$, is a mapping from the output places of t to vacant cells in the current marking.*

Since each token passes through a net level transition when it fires, a transition is enabled only when there is a bijective function between the chosen enabling and an output.

DEFINITION 4 (CONCESSION). *A transition t has* concession *for a firing of type α if there is an enabling* $\mathrm{E}(t, \alpha)$

[1]We assume all arcs have multiplicity 1.

such that there is a bijective mapping ϕ from $\mathrm{E}(t,\ \alpha)$ *to an output* $O(t)$*, which preserves the types of tokens.*

As with classical Petri nets with priority, having concession identifies those transitions which could legally fire according to the net structure and the current marking. Those transitions which *can* fire are determined by the priorities.

DEFINITION 5 (ENABLING RULE). *A transition t will be* enabled *for a firing of type α if there is no other transition of higher priority with concession in the current marking.*

DEFINITION 6 (FIRING RULE). *When a transition t fires with type α on the basis of the enabling* $\mathrm{E}(t,\alpha)$*, and concession ϕ then for each* $(\mathbf{P}_i,\ T,)$ *in* $\mathrm{E}(t,\ \alpha)$*, $T[T]$ is replaced by $T[_]$ in the marking of $\mathbf{P}_i$, and the current marking of each output place is updated according to ϕ.*

We assume that when there is more than one mapping ϕ from an enabling to an output, then they have equal probability and one is selected randomly. The rate of the enabled firing is determined using apparent rates, and the notion of bounded capacity, as usual for PEPA.

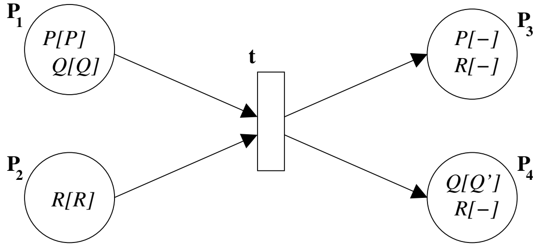


Figure 2: PEPA net fragment to illustrate enabling

To illustrate the firing rules we consider the small PEPA net fragment shown in Figure 2. Here we assume that transition **t** has type α and that each of the tokens P, Q, and R has a one-step α-derivative (P', Q' and R' respectively). For the net fragment shown there are two enablings:

$$((\mathbf{P_1}, P), (\mathbf{P_2}, R)) \quad \text{and} \quad ((\mathbf{P_1}, Q), (\mathbf{P_2}, R))$$

There are also two outputs:

$$((\mathbf{P_3}, P[_]), (\mathbf{P_4}, R[_])) \quad \text{and} \quad ((\mathbf{P_3}, R[_]), (\mathbf{P_4}, R[_]))$$

Note that $((\mathbf{P_3}, P[_]), (\mathbf{P_3}, R[_]))$ is not an output.

Here t has concession because there is a bijective mapping ϕ from $((\mathbf{P_1}, P), (\mathbf{P_2}, R))$ to $((\mathbf{P_3}, P[_]), (\mathbf{P_4}, R[_]))$. Note that the other potential mappings,

$$((\mathbf{P_1}, Q), (\mathbf{P_2}, R)) \quad \text{to} \quad ((\mathbf{P_3}, P[_]), (\mathbf{P_4}, R[_]))\,,$$
$$((\mathbf{P_1}, P), (\mathbf{P_2}, R)) \quad \text{to} \quad ((\mathbf{P_3}, R[_]), (\mathbf{P_4}, R[_]))\,, \text{and}$$
$$((\mathbf{P_1}, Q), (\mathbf{P_2}, R)) \quad \text{to} \quad ((\mathbf{P_3}, R[_]), (\mathbf{P_4}, R[_]))$$

are not valid because the types of tokens would not be preserved.

If no other transitions of higher priority were enabled in the PEPA net this transition would fire, removing a P token from place P_1, and a R token from place $\mathbf{P_2}$ and placing a P' token in place $\mathbf{P_3}$ and a R' token in place $\mathbf{P_4}$.

3. TRANSLATING PEPA NETS TO PEPA

In order to be able to exploit fully the existing suite of analysis tools for PEPA models we have defined a translation from PEPA nets to PEPA. In this section we describe the algorithm to do this in some detail. The algorithm is composed of a number of different steps which are described in turn in the following sections.

Before starting the translation, a *preprocessing* phase in necessary in order to rename firings that in the PEPA nets share the same input place and same action type. This is done by checking the arcs specification in the PEPA net file. For example, we need to make transformations of the arcs such as the following

$$\mathbf{P}_k \xrightarrow{(\alpha,r_1)} \mathbf{P}_j \quad \text{and} \quad \mathbf{P}_k \xrightarrow{(\alpha,r_2)} \mathbf{P}_i$$
$$\rightsquigarrow$$
$$\mathbf{P}_k \xrightarrow{(\alpha_1,r_1)} \mathbf{P}_j \quad \text{and} \quad \mathbf{P}_k \xrightarrow{(\alpha_2,r_2)} \mathbf{P}_i$$

and to store the triples (α, α_1, r_1) and (α, α_2, r_2) in an array denoted by *Fire*.

Several data types and functions are used in the algorithm and are introduced below. In addition we use $\vec{\mathcal{A}}(C)$, returning the complete action type set of component C.

All	stores all action types of the PEPA model
NumberSC[ ]	counts occurrences of static components
Numberα_j	counts renamings of activity α_j as introduced by preprocessing
OccurSC[ ]	stores repeated static components
Fire[ ]	stores the triples from the preprocessing
TargetP$_j$	stores triples (P, α_j, r) where P is the output place of the arc and (α_j, r) its label
Collect[α]	stores the components which must synchronise on the current action type α
Extract(α_j)	returns the original action type α from the triple in *Fire*; *Nil* otherwise
$\mathcal{L}_{sync}(T_i, SC_k)$	synchronisation set between token type T_i and the static component SC_k
$\mathcal{T}(C)$	returns the type of component C

3.1 Step 1: Translating Static Components

This step concerns the translation of static components. We need to count the occurrences of static components within the places of the PEPA net and rename action types and derivatives of each occurrence of the static component to reflect its location. This allows us to avoid erroneous synchronisations since in a PEPA net synchronisation is possible only when the components are resident in the same place.

for each static component SC_i

> $\mathit{NumberSC}[i] \leftarrow 0$;
> $\mathit{OccurSC}[i] \leftarrow \emptyset$;
> **for each** place $P_j \in \mathcal{P}$
>   **if** (SC_i in P_j)
>   **then** $\mathit{NumberSC}[i] \leftarrow \mathit{NumberSC}[i] + 1$;
> **if** ($\mathit{NumberSC}[i] > 1$) **then**
> $\mathit{OccurSC}[i] \leftarrow SC_i$;
>   **for** $k = 1 \ldots \mathit{NumberSC}[i]$
>   make an instance SC_{ik} of SC_i, renumbering activities and derivatives;
>   *// save all action types in All*
> $\mathit{All} \leftarrow \mathit{All} \cup \vec{\mathcal{A}}(SC_{ik})$;

3.2 Step 2: Translating Cells

In this step we consider all the PEPA net places containing cells and we associate one component with each cell within a place. Each cell definition is built by considering the input and output arcs to the place itself, with respect to the type of the token which may occupy the cell. Moreover, each cell can be in two states: *empty*, denoted by a subscript 0 in its derivative, and *full* denoted by a subscript 1. Firing a token into the place changes the state of the cell from empty to full. Firing a token out of the cell changes its state from full to empty. This allows us to prevent a token from moving to a place where all of the cells are already full.

for each token type T and place $P_j \in \mathcal{P}$

> **for each** cell i in P_j
>   create a component Cell_{Tji0};
>   **for each** input arc $P_k \xrightarrow{(\alpha,r)} P_j$ such that $\alpha \in \vec{\mathcal{A}}(()T)$
> $\mathit{All} \leftarrow \mathit{All} \cup \{\alpha\}$;
>   *// avoiding redundant activities*
> $\mathit{Cell}_{Tji0} \leftarrow \mathit{Cell}_{Tji0} + (\alpha, r).\mathit{Cell}_{Tji1}$;
>   **for each** output arc $P_j \xrightarrow{(\alpha,r)} P_k$ such that $\alpha \in \vec{\mathcal{A}}(()T)$
> $\mathit{All} \leftarrow \mathit{All} \cup \{\alpha\}$;
> $\mathit{Cell}_{Tji1} \leftarrow \mathit{Cell}_{Tji1} + (\alpha, r).\mathit{Cell}_{Tji0}$;

3.3 Step 3: Translating Tokens

Token definitions need to be revised according to their dynamic behaviour. Two aspects are taken into account: the movement of a token into a new place after a firing and its interaction with static components within a PEPA net place. Again, in order to avoid erroneous synchronisations, we need to create new names for action types and derivatives. This is done when two or more net transitions have the same label and when tokens synchronise with static components which appear several times in the net. Recall that in Step 1 the action types and the derivatives of each occurrence of the static component were renamed. Therefore here we need to copy the derivatives of the token where the synchronising action type appears and we do this for each renaming of the synchronising action.

for each token type T_i

> *// movement of the token in the net*
> **for each** action type $\alpha_j \in \mathcal{A}_f(T_i)$
> $\mathit{Number}\alpha_j \leftarrow 0$;
> $\mathit{TargetP}_j \leftarrow \emptyset$;
>   **for each** arc $P_k \xrightarrow{(action,r)} P_l$
> $\alpha_j^* \leftarrow \mathit{Extract}(action)$;
>   **if** ($\alpha_j = \alpha_j^*$) **then**
> $\mathit{Number}\alpha_j \leftarrow \mathit{Number}\alpha_j + 1$;
> $\mathit{TargetP}_j \leftarrow \mathit{TargetP}_j \cup P_l$;
>   **if** ($\mathit{Number}\alpha_j > 1$) **then**
> $r^* \leftarrow r/\mathit{Number}\alpha_j$;
>   **for each** distinct place in $\mathit{TargetP}_j$
>   add the corresponding activity $(action, r^*)$ in $\mathit{TargetP}_j$ followed by a new component identifier C^*;
>   duplicate the sequential component C resulting from the execution of action α_j in T_i, replacing C with C^*;
>   **else** make a copy of the current derivative of T_i;
> *// interaction of token and static components*
> **for each** place $P_l \in \mathcal{P}$
>   **if** (T_i in P_l) **then**
>   **for each** static component SC_k in $\mathit{OccurSC}$
>   **for each** action type $\alpha_j \in \mathcal{L}_{sync}(T_i, SC_k)$
>   **for** $x = 1 \ldots \mathit{NumberSC}[k]$
>   duplicate the derivatives where α_j appears in T_i for each renaming of α_j during the translation of the static component SC_k;
> *// representing the initial marking*
> **for each** occurrence of T_i at P_j in M_0
>   make an instance ST_{ij} of T_i, renumbering activities and derivatives.

3.4 Step 4: Building the System Equation

The "system equation" of a PEPA model forms a parallel composition of the process components which have defining equations. Several copies of each component might be required and any instance of a process definition can be configured by the use of a synchronisation set which forces it to synchronise with one or more of the other processes. We build the system equation for the PEPA model generated from the input PEPA net by putting in parallel all the components generated in the previous steps. Then we force them to synchronise on common action types.

// *Build the equation with empty synchronising sets*
take the first component C_0 in the PEPA model

$System \leftarrow\ '(' + C_0;$

for each component C_i

$System \leftarrow System\ +\ '(<>'\ + C_i$

balance the number of opened left parenthesis

// *Fill the synchronising sets in the system equation*

remove repeated action types from *All*

for each $\alpha \in All$

for each component C

if $\alpha \in \mathcal{A}(C)$ **then**

$Collect[\alpha] \leftarrow C;$

for each pair $(C_k, C_l) \in Collect[\alpha]$

if $\mathcal{T}(C_k) \neq \mathcal{T}(C_l)$ **then**

insert α in the synchronisation set associated with the most left component between C_k and C_l

3.5 Properties of the translation

PEPA stochastic process algebra models have no notion of context, nor any capacity to move components from one context to another, and therefore cannot express the concept of dynamically varying communication structure. The task performed by the compilation of a PEPA net described above has been to remove all of the mobility from the PEPA net by making components' behaviour depend on location. This is achieved by expanding the definition of the tokens of the net replicating local state behaviours and customising these for each cell which the token may visit.

The compilation from a PEPA net to a PEPA model is analogous to the preprocessing of a coloured Petri net to transform it into an uncoloured net for analysis purposes. In the worst case—when there is only a single class of token, all tokens can travel to all of the places of the net, and the net has no static components—the resulting PEPA model is n times larger than the input PEPA net, where n is the number of cells in the net. Thus there is no exponential increase in the size of the representation of the model here and the compilation procedure runs in a time proportional to the size of the model, not the size of the state space which it would generate when fired against the operational semantics of the language.

The correctness of the translation from the input PEPA net to the resulting PEPA model can be stated formally thus. When evaluated against the operational semantics of the languages the two models generate labelled transition systems (LTS) for the models. These LTS contain information about the activities performed in the model by including action types in labels. If these occurrences of action types are removed then isomorphic CTMCs are obtained for the input PEPA net and its translation. This is the sense in which the output PEPA model is a correct translation of the input PEPA net.

4. EXAMPLE: MOBILE AGENTS

We present a small example to reinforce the reader's understanding of the translation algorithm. In this example a mobile software agent visits three sites. It interacts with static software components at these sites and has two kinds of interactions. When visiting a site where a network probe is present it interrogates the probe for the data which it has gathered on recent patterns of network traffic. When it returns to the central co-ordinating site it dumps the data which it has harvested to the master probe. The master probe analyses the data.

4.1 The PEPA net Model

The structure of the application is as represented by the PEPA net in Figure 3. This marking of the net shows the mobile agent resident at the central co-ordinating site. Recall that as a memory aid, in both the net and the PEPA token definitions, we print in bold the names of those activities which can cause a firing of the net. In this example the activities which can cause a firing of the net are **go** and **return**.

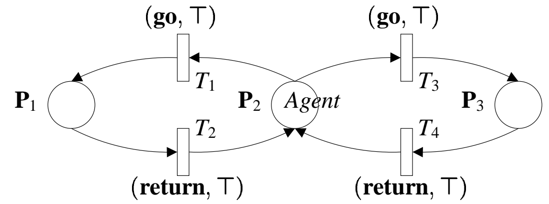


Figure 3: A simple mobile agent system

Formally, we define the places of the net as shown in the PEPA context definitions below.

$$\mathbf{P}_1[Agent] \stackrel{def}{=} Agent[Agent] \underset{\{interrogate\}}{\bowtie} Probe$$
$$\mathbf{P}_2[Agent] \stackrel{def}{=} Agent[Agent] \underset{\{dump\}}{\bowtie} Master$$
$$\mathbf{P}_3[Agent] \stackrel{def}{=} Agent[Agent] \underset{\{interrogate\}}{\bowtie} Probe$$

The behaviour of the components is given by the following PEPA definitions.

$$\begin{aligned}
Agent &\stackrel{def}{=} (\mathbf{go}, \lambda).Agent' \\
Agent' &\stackrel{def}{=} (interrogate, r_i).Agent'' \\
Agent'' &\stackrel{def}{=} (\mathbf{return}, \mu).Agent''' \\
Agent''' &\stackrel{def}{=} (dump, r_d).Agent \\
Master &\stackrel{def}{=} (dump, \top).Master' \\
Master' &\stackrel{def}{=} (analyse, r_a).Master \\
Probe &\stackrel{def}{=} (monitor, r_m).Probe + \\
&\quad\ (interrogate, \top).Probe
\end{aligned}$$

The initial marking of the net has one *Agent* token in place $\mathbf{P}_2$ and no other tokens:

$$\begin{array}{llll}
(& Agent[_] & \underset{\{interrogate\}}{\bowtie} & Probe, \\
& Agent[Agent] & \underset{\{dump\}}{\bowtie} & Master, \\
& Agent[_] & \underset{\{interrogate\}}{\bowtie} & Probe \quad)
\end{array}$$

4.2 Translation into PEPA

The example has two static components *Master* and *Probe*. As the latter appears twice in the PEPA net model (in places $\mathbf{P}_1$ and $\mathbf{P}_3$), the application of the algorithm generates two instantiations of this component $Probe_1$and $Probe_2$. Thus in the PEPA model we have:

$$
\begin{array}{lcl}
Master & \stackrel{def}{=} & (dump, \top).Master' \\
Master' & \stackrel{def}{=} & (analyse, r_a).Master \\
 & & \\
Probe_1 & \stackrel{def}{=} & (monitor_1, r_m).Probe_1 \\
 & + & (interrogate_1, \top).Probe_1 \\
 & & \\
Probe_2 & \stackrel{def}{=} & (monitor_2, r_m).Probe_2 \\
 & + & (interrogate_2, \top).Probe_2
\end{array}
$$

As each place contains just one cell, the translation algorithm will generate three new components which we will denote by $Cell_{j0}$, where the subscript $j = 1$, 2 or 3 specifies that the cell resides at place $\mathbf{P}_1$, $\mathbf{P}_2$ or $\mathbf{P}_3$. The suffix which follows j takes only the values zero or one and is used to record whether the cell is empty or full[2].

Firings and transitions are no longer distinguished because there is only one class of activities in PEPA and so we no longer embolden the names *go* and *return*.

$$
\begin{array}{lcl}
Cell_{10} & \stackrel{def}{=} & (go_1, \top).Cell_{11} \\
Cell_{11} & \stackrel{def}{=} & (return_1, \top).Cell_{10} \\
 & & \\
Cell_{20} & \stackrel{def}{=} & (return_1, \top).Cell_{21} + (return_2, \top).Cell_{21} \\
Cell_{21} & \stackrel{def}{=} & (go_1, \top).Cell_{20} + (go_2, \top).Cell_{20} \\
 & & \\
Cell_{30} & \stackrel{def}{=} & (go_2, \top).Cell_{31} \\
Cell_{31} & \stackrel{def}{=} & (return_2, \top).Cell_{30}
\end{array}
$$

Component *Agent* is the token in the net and its definition is modified to take into account its complete dynamic behaviour in the net.

$$
\begin{array}{lcl}
Agent & \stackrel{def}{=} & (go_1, \lambda/2).Agent' + (go_2, \lambda/2).Agent_1' \\
Agent' & \stackrel{def}{=} & (interrogate_1, r_i).Agent'' \\
Agent_1' & \stackrel{def}{=} & (interrogate_2, r_i).Agent_1'' \\
Agent'' & \stackrel{def}{=} & (return_1, \mu).Agent''' \\
Agent_1'' & \stackrel{def}{=} & (return_2, \mu).Agent''' \\
Agent''' & \stackrel{def}{=} & (dump, r_d).Agent
\end{array}
$$

The PEPA system equation forms a parallel composition of instances of the above components and requires them to synchronise on their common actions. The system equation in this case is as follows:

$$
\begin{array}{lcl}
System & \stackrel{def}{=} & (Cell_{10} \underset{K_1}{\bowtie} (Probe \underset{K_2}{\bowtie} (Agent \underset{K_3}{\bowtie} (Cell_{21} \\
 & & \underset{K_4}{\bowtie} (Master \underset{K_5}{\bowtie} (Probe_1 \underset{K_6}{\bowtie} Cell_{30}))))))
\end{array}
$$

where $K_1 = \{ go_1, return_1 \}$, $K_2 = \{ interrogate_1 \}$, $K_3 = \{ go_1, return_1, go_2, return_2, dump, interrogate_2 \}$, $K_4 = \{ go_2 \}$, $K_5 = K_6 = \emptyset$.

[2]Note that since there is only one token type in this PEPA net we omit the corresponding subscript.

5. IMPLEMENTATION

We have implemented our translation from PEPA nets to PEPA as an extension to our existing tool, The PEPA Workbench for PEPA Nets, ML Edition. This is a modular research version of the PEPA Workbench for experimentation and extension so it is well-suited to adding modules such as the PEPA net compiler.

The workbench front end which consists of the PEPA net parser and syntax tree builder required no changes from their use as components of the PEPA Workbench. Similarly, the static analyser for PEPA nets was unchanged. The new components added by the compiler make use of this front end and then provide other functions for traversing the abstract syntax tree of the PEPA net and inferring derived synchronisation sets from the union of sets which are stored throughout the tree. Similarly, token flows were inferred for the components which flow around the net. Additionally, new names were coined for specialisations of activities within cells. Finally, a PEPA "system equation" is generated which composes the images of tokens, cells and static components using the appropriate inferred synchronisation sets.

We have found the implementation of the compiler to be very efficient in practice. The PEPA net compiler does not unfold the global state space of the model. Instead it uses the compositional nature of the PEPA net when constructing an equivalent PEPA model. In this way, the compilation is performed in linear time, with respect to the size of the original PEPA net description.

6. PLATFORM SUPPORT FOR PEPA

As mentioned earlier, the objective of the translation from the PEPA nets modelling language into the PEPA stochastic process algebra is to allow the PEPA net modeller to exploit the existing suite of tools which support PEPA. In this section we give a brief overview of these tools. We will use these tools in processing our case study example in Section 7 so we provide an explanation of these tools here before we use them later.

We see the provision of tool support for modelling languages as being analogous to providing tools for programming languages. A software development kit typically provides a range of tools (compilers, debuggers, profilers, perhaps even model checkers) which perform various types of analysis or conversion on the program. Similarly in providing platform support for PEPA and PEPA nets we provide this through a modelling kit containing solvers, passage-time analysers, model-checkers and other tools.

The analysis of a PEPA model proceeds by deriving its underlying Markov chain and solving this to find the long-run probability of the states of the chain. The states of the chain are in one-to-one correspondence with the states of the derivation graph of a PEPA process as specified by the operational semantics of the language so information about the long-run behaviour of the CTMC translates directly to information about the PEPA model from which it was derived.

The PEPA Workbench [3] generates the CTMC corresponding to any given PEPA model. The Workbench is available in

two editions (principally implemented in ML and Java respectively), connected to a variety of solvers for the CTMC. Additionally, in 1999 Clark implemented some additional PEPA tools such as the PEPAroni simulator in Pizza [4], an extension of Java. This was later re-implemented in the Java version of the Workbench [5].

The PEPA Workbench is available for download from the PEPA Web site located at `http://www.dcs.ed.ac.uk/pepa/` together with documentation, example models and copies of papers on PEPA and PEPA nets.

6.1 Analysis Tools

Generating the underlying CTMC, and finding its steady state probability vector is rarely, if ever, the objective of PEPA or PEPA net modelling. Formal tool support for querying performance models is an area which has received little attention until recently, despite its practical importance.

At the most basic level the modeller wishes to construct a *reward structure* over the state space of the CTMC, to be used in conjunction with the steady state probability vector to derive performance measures. For steady state measures the reward structure is a vector recording a reward for each state, although for many states the reward value will be zero. Thus the problem becomes one of identifying the appropriate set of states to attach a non-zero reward to. Clearly, when the CTMC arises from a stochastic process algebra model we prefer to characterise the state at the process algebra level. PEPA analysis tools have been developed which tackle this problem in two distinct ways.

6.1.1 The PEPA State Finder

The PEPA State Finder identifies subsets of states using regular expression pattern matching, applied to the table of PEPA expressions which make up the state of the model. Recall that there is a one-to-one correspondence between the syntactic forms of the PEPA process as it evolves and the states of the CTMC. The Workbench maintains a table recording this correspondence, and using regular expression pattern matching the PEPA State Finder is able to extract the states of interest. For example, it is possible to use an expression such `*|(next,r).*` to return the state numbers of all the states in which the second component enables a (*next*, *r*) activity. This could then be used to construct a reward structure suitable for calculating the throughput of *next* in the second component: the value *r* is placed in the reward vector at each position corresponding to a (numerical) state found by the PEPA State Finder.

6.1.2 PML_μ

A more sophisticated means of specifying rewards is described in Clark's PhD thesis [6], and developed around the stochastic logic PML_μ. Inspired by the probabilistic model logic of Laren and Skou, PML [7], PML_μ is able to differentiate PEPA terms which perform the same activities but at different rates. The key to this is a modification to Hennessy-Milner logic in which the diamond operator becomes decorated with a rate. The semantics of an expression in the logic is a subset of states, and thus logical expressions may be used, in conjunction with a value, to specify a reward structure. Clark extended the ML edition of the PEPA Workbench to include support for PML_μ and associated reward structures [6].

6.1.3 CSL

More recently the CSL logic (Continuous Stochastic Logic) [8] has gained some acceptance as a suitable vehicle for expressing performance and performability measures which can be model checked on a CTMC. Unlike the approach of PML_μ, where the logic formula is used to distinguish a subset of states (those satisfying the formula), a CSL formula expresses an assertion about the performance measures of a model which can then be checked to see whether it is true or not. Thus this offers an alternative approach to using logics to assist in performance analysis.

6.2 Interoperation with other tools

In addition to developing tools specific to PEPA, when possible we connect the PEPA language to existing tools for performance modelling.

6.2.1 PEPA and Möbius

The Möbius modelling framework [9] is a multi-formalism and multi-paradigm modelling tool, and as such provided an excellent tool to integrate PEPA into [10]. Integration in this case involved not only building PEPA support into Möbius but also incorporating Möbius features into the version of PEPA which was implemented in Möbius. The reason for this was that it then became possible to share variables between components modelled in PEPA and components implemented in another Möbius formalism, such as SANs.

Building support for PEPA directly in another tool has the advantage that the language is supported efficiently, without additional overheads imposed by translation from one representation into another. However, it has the implementation cost that the implementor must be familiar not only with the concepts of the host tool but also with their representation in data structures and algorithms. Such tight integration is not always cost-effective or necessarily desirable.

6.2.2 PEPA and PRISM

For our subsequent integration work we used a lighter-weight, component-based method involving a relatively loose coupling between the language and the host solver. This was adopted when a binding for the PEPA language in the PRISM [11] probabilistic model checker was developed. In PRISM analysis of CTMCs is performance through model checking specifications in the probabilistic temporal logic CSL.

In order to define and analyse a model, PRISM requires two input files: a *description of the system* under investigation and a *set of properties* to be checked against it. The description of the system is given in a reactive module language. To integrate PEPA, a PEPA compiler was adapted to generate its output in the form of this language. Additionally, the model checker itself was extended to support PEPA's combinators (parallel and hiding).

This method of working with PEPA models requires a significant degree of expertise on the part of the modeller, because errors in evaluation can occur in the tool chain right from the PEPA parser through PEPA-to-PRISM and PRISM down to CUDD [12], the BDD library which provides MTBDD data structures and algorithms to PRISM. Even if errors do not occur in translation still to achieve the best performance from the solver it is necessary to know how to configure both PRISM and CUDD which means that this method is best suited to experienced modellers only.

The PRISM tool is available for download from `www.cs.bham.ac.uk/~dxp/prism/`.

6.2.3 PEPA, IPC and Dnamaca

A recent development in the PEPA tools is Bradley's ipc (The Imperial PEPA Compiler) [13]. The ipc tool translates an input PEPA model into the Petri net notation provided by Knottenbelt's Dnamaca tool [14]. The ipc tool supports the PEPA language comprehensively. Apparent rates are supported, as are anonymous components. These are two advantages over the PEPA-to-PRISM compiler, and a richer class of PEPA models can therefore be analysed by ipc/Dnamaca as a result.

In one other important respect, ipc provides more comprehensive PEPA support than comparable tools because it also translates PML$_\mu$ formulae into the Dnamaca specification language. The Dnamaca specification language is a classical Petri net logic allowing specification formulae to quantify the number of tokens in the places of the net and thereby identify states and sets of states within the reachable state space of the model.

Via ipc, the unique solution capabilities of Dnamaca become available and because of this it is possible to perform efficient passage time analysis over PEPA models.

The ipc compiler is available for download from `www.doc.ic.ac.uk/ipc`.

7. EXAMPLE: SECURE WEB SERVICE

In this section we present a PEPA net model of an example taken from the field of Web Services, a relatively new Web programming paradigm which provides standard interfaces and communication protocols for integrating Web-based distributed applications.

Web services allow different applications from different sources to communicate and interoperate. Many practitioners believe that they will become a key component of many organization's future business-integration initiatives. Security issues need to be considered carefully because the integration of applications that communicate over a network could reveal sensitive information to unknown parties.

Web services interfaces are defined, described and discovered by means of XML—the eXtensible Markup Language—and therefore they are not tied to any single operating system or programming language. XML is used to tag the data which are exchanged between applications via SOAP (Simple Object Access Protocol) messages, XML-based messages which may be transported using a variety of Internet protocols, including SMTP, MIME, and HTTP.

The available services are described using WSDL (the Web Services Description Language), an XML-formatted language, and they are made available thanks to UDDI (the Universal Description, Discovery and Integration format), a Web-based distributed directory that lists the available services, similar to a traditional phone book's yellow and white pages.

Our example is a model of a mobile object system where a client sends SOAP message objects to a remote Web service. Our scenario is that a financial tycoon is sending requests for stocks and share price information to a remote Web service which provides this information. These requests for information are encrypted. An eavesdropper could make use of the information in the messages if they were sent as clear text. The Web service itself is protected by a firewall.

7.1 The token type

The tokens exchanged in the system are SOAP messages in various formats. These may either be sent across the network as clear text or encrypted to preserve their contents. A SOAP message may be parsed to build an in-memory data structure which can be read and modified as needed. This data structure is a DOM tree (a Document Object Model tree).

$$\begin{array}{rcl} \mathit{SoapMessage} & \stackrel{\mathit{def}}{=} & (\mathit{send}_{\mathit{clr}}, r_{\mathit{sc}}).\mathit{SentClearMessage} \\ & + & (\mathit{encrypt}, r_e).\mathit{EncryptedMsg} \\ & + & (\mathit{parse}, r_p).\mathit{DOMtree} \\ \mathit{SentClearMessage} & \stackrel{\mathit{def}}{=} & (\mathbf{copyClear}, \top).\mathit{SoapMessage} \end{array}$$

The rate at which information is copied across the network is specified not in this component but at the upper PEPA net level. The description of this component leaves unspecified ($\top$) the rates at which these actions are performed, allowing the cooperating partner in the synchronisation to determine these rates. All such unspecified rates must be synchronised with specified ones in the final model. Frequently, synchronisations are between one active partner and one passive partner.

Encrypted messages can be decrypted to recover their initial contents or sent across the network in encrypted form.

$$\begin{array}{rcl} \mathit{EncryptedMsg} & \stackrel{\mathit{def}}{=} & (\mathit{decrypt}, r_d).\mathit{SoapMessage} \\ & + & (\mathit{send}_{\mathit{enc}}, r_{\mathit{se}}).\mathit{SentEncMessage} \\ \mathit{SentEncMessage} & \stackrel{\mathit{def}}{=} & (\mathbf{copyEncrypted}, \top).\mathit{EncryptedMsg} \end{array}$$

In both cases, we model the transmission of a SOAP message as a two-phase process, separating the cost of making the decision to send ($\mathit{send}_{\mathit{clr}}$, $\mathit{send}_{\mathit{enc}}$) from the cost of copying the bytes across the network (**copyClear**, **copyEncrypted**).

DOM trees may be read or modified. As an in-memory data structure they first must be serialised (using the *export* activity) if they are to be sent across the network.

$$\begin{array}{rcl} \mathit{DOMtree} & \stackrel{\mathit{def}}{=} & (\mathit{read}, r_r).\mathit{DOMtree} \\ & + & (\mathit{modify}, r_m).\mathit{DOMtree} \\ & + & (\mathit{export}, r_x).\mathit{SoapMessage} \end{array}$$

7.2 Static components

SOAP message tokens of various forms are exchanged between the places of the net. Static components at these places interact with the tokens. On the client side is the user, making requests of the remote Web Service. The user encrypts requests before they are sent and decrypts replies when they are received.

$$\begin{aligned} User &\stackrel{def}{=} Encrypt + Decrypt \\ Encrypt &\stackrel{def}{=} (encrypt, \top).(send_{enc}, \top).User \\ Decrypt &\stackrel{def}{=} (decrypt, \top).(parse, \top).(read, \top).Request \\ Request &\stackrel{def}{=} (modify, \top).(export, \top).User \end{aligned}$$

Running on the firewall is a gatekeeper process which performs three distinct functions: decrypting user requests, bouncing flawed requests and encrypting replies from the server. The gatekeeper receives requests from the user and decrypts them. The decrypted message might be a well-formed request, in which case it is forwarded to the server. Alternatively it might be in an invalid format, request a non-existent service, or have suspicious attachments. In this case it is bounced back to the user with a diagnostic error message attached. This assignment of responsibilities means that the load on the server is reduced.

All of the communication with the user is sent in encrypted format. Behind the firewall the messages are exchanged in the clear. Thus $send_{enc}$ always sends to the user and $send_{clr}$ always sends to the server.

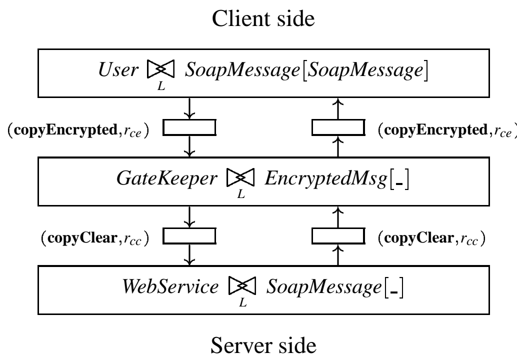
$$\begin{aligned} GateKeeper &\stackrel{def}{=} FilterIn + Bounce + FilterOut \\ FilterIn &\stackrel{def}{=} (decrypt, \top).(send_{clr}, \top).GateKeeper \\ Bounce &\stackrel{def}{=} (decrypt, \top).FilterOut \\ FilterOut &\stackrel{def}{=} (encrypt, \top).(send_{enc}, \top).GateKeeper \end{aligned}$$

Behind the firewall the share price Web service runs on the server. Its life cycle is parsing, reading, modifying, serialising and returning requests.

$$\begin{aligned} WebService \stackrel{def}{=} & (parse, \top).(read, \top). \\ & (modify, \top).(export, \top).(send_{clr}, \top).WebService \end{aligned}$$

7.3 The PEPA net

The PEPA net of the system sites the above static components at places of the net and specifies the communication between different places of the net by naming the transitions which must be fired for tokens to move from place to place. The *User* is operating the client machine, the *GateKeeper* process runs on the firewall, the *WebService* on the server behind the firewall.



The synchronisation set used at each place, L, is $\{$ *decrypt*, $send_{clr}$, *parse*, *read*, *modify*, *export*, *encrypt*, $send_{enc}$ $\}$.

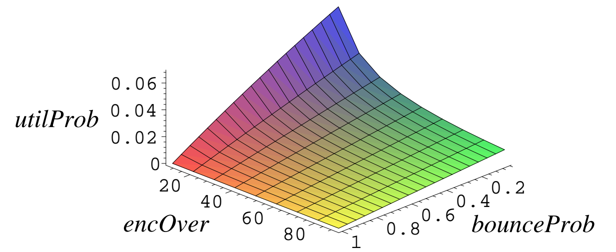
7.4 Model analysis

We analysed our secure Web service model using the PEPA Workbench for PEPA nets and the PRISM model checker, making use of our PEPA net compiler tool to compile the net into a PEPA model which is accepted by PRISM.

7.4.1 Using the PEPA Workbench for PEPA nets

The PEPA Workbench for PEPA nets implements the semantics of the PEPA nets language directly and generates a CTMC suitable for performance analysis.

In the analysis of our example using the PEPA Workbench we calculated the utilisation of the Web service as two quantities varied. The first was the probability (*bounceProb*) that user requests to use the Web service would be bounced by the gatekeeper process (for whatever reason). We varied this between 0.05 (there is a low probability that requests will be bounced) and 1.0 (requests messages are always bounced, and do not reach the server). The second parameter which we varied was the strength of the encryption used to encrypt messages. This variation is captured by an increase in the overhead associated with encryption and decryption (*encOver*). We plotted this for strong encryption taking in the range ten times longer to compute to ninety times longer.



As expected, the utilisation of the Web service is seen to be lowest when the strongest form of encryption is used (because this delays only the user and the gatekeeper) and when the probability of bouncing messages is highest (in that case they never reach the server).

7.4.2 Using the PEPA nets compiler

The PEPA nets Compiler is a component of the PEPA Workbench for PEPA nets. It compiles the secure Web service example from a PEPA net into an equivalent PEPA process algebra model in 0.02 seconds on a 1.6GHz Pentium IV processor machine with 256Mb of memory. The GNU `time` command version 1.7 was used to obtain this timing information. The time reported is elapsed real (wall clock) time used by the process. This compilation facilitates further analysis with tools which support PEPA.

7.4.3 Model-checking using PRISM

We used PRISM to model check a range of CSL formulae over the PEPA model generated by the PEPA net compiler. One formula which we checked was “with probability at least 0.5,

the web service begins processing its first request within t time units". This is expressed in CSL as follows:

$$\mathrm{P}_{\geq 0.5}[\textbf{true}\ \mathrm{U}_{\leq t}(\mathit{WebService} = (\mathit{read}, \top).\mathit{WebService}')].$$

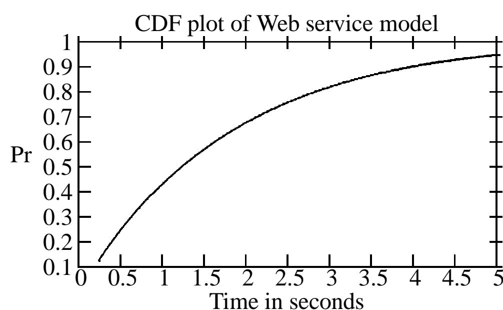
This is a typical CSL formula, combining time, probability and a path through the system evolution. We checked this formula with the time parameter, t, taking values between 1.0 and 100,000.0 time units. Model-checking formulae of this type on a 1.6GHz Pentium IV processor machine with 256Mb of memory typically requires an average of 0.275 seconds of compute time, as reported by PRISM. These times are measured in the PRISM implementation using Java's native method System.currentTimeMillis(). The formula above is satisfied for all states of the secure Web Service model for values of t=250 time units and higher.

7.4.4 Passage-time analysis using ipc/DNAmaca

Some questions which we would like to ask of a performance model are not readily answered from the computation of the steady-state probability distribution of the model, nor are they answered by model-checking time-bounded probabilistic formulae against the model. Queries such as those involving responsiveness of the system under workload—"does the system respond within 5 seconds 90% of the time?"—are instead answered by reference to passage-time quantiles.

To undertake analysis of this form we use the Imperial PEPA Compiler (ipcv0.91a 230603, James Clerk Maxwell Release) together with the DNAmaca analyser for Markov and semi-Markov chains (Hydra release). This tool chain has previously been used to analyse PEPA models [15] but this is the first application of these tools to a PEPA net.

The plot below shows the cumulative density function for the passage from the user initiating a web service request on the client side to the initiated computation being completed at the server side.



In order to effect the passage-time analysis of this model we used the method of *stochastic probes* [16] whereby additional stochastic process algebra components are generated to probe the system to observe the specified start events and end events. The resulting extended PEPA model is then compiled by ipc into a DNAmaca model.

Specifically in this case the PEPA actions which we probed were send_{enc} at the *Client* place and send_{clr} at the *Server* place in the net. The PEPA net compiler distinguishes these in the resulting PEPA stochastic process algebra model so that we can be sure that we are probing these activities and not the send_{enc} and send_{clr} activities which occur at the *Firewall* place in the PEPA net.

8. CONCLUSIONS

Building performance models of realistic real-world systems is an activity which requires careful attention to detail in order to model correctly the intended behaviour of the system. Proceeding with care during this part of the modelling process is a wise investment of effort. If the initial performance model contains errors then all of the computational expense incurred in solving the model and all of the intellectual effort invested in the analysis and interpretation of the results obtained would at best be wasted. In general interpreting a model with errors could lead to making flawed economic or strategic decisions based on erroneous conclusions made from erroneous results. For this reason we consider it important to work with structured, high-level modelling languages which directly support the concepts and idioms of the application domain, such as code mobility. In this paper we have shown how a mobile code modelling language, PEPA nets, can be mapped into another language, the stochastic process algebra PEPA, where mobility is not directly expressible in the language and there is no notion of physically or logically separated locations being a barrier to communication.

Our translation from PEPA nets to PEPA has enough subtleties that it was worthwhile to implement it in a software tool to prevent making mistakes that could easily be made if applying the method by hand. We extended our existing workbench for PEPA nets for this purpose, building in a compiler from PEPA nets to PEPA. We applied this to a case study of modelling a secure Web Service. We analysed the resulting performance model for steady-state behaviour, transient behaviour exposed by CSL model-checking, and responsiveness quantified by the computation of passage-time quantiles.

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